

### Homework #5

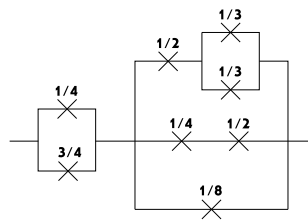
Due Thursday, June 7, 2018, at 2:30 PM

*Collaboration and discussions are not allowed on Problem 1  
and are allowed and encouraged on Problem 2*

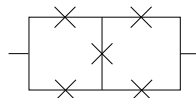
1. **Stochastic Switching Circuits** (*Collaboration is not allowed on this Problem*)

In this problem we study the synthesis and analysis of stochastic switching circuits. In particular, we consider simple series-parallel circuits, series-parallel circuits and general circuits.

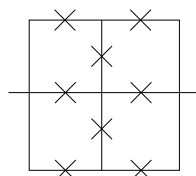
- (a) In class we discussed an algorithm (called the B-algorithm) for synthesizing simple series-parallel circuits for realizing probabilities of the form  $a/2^n$ , where  $0 \leq a \leq 2^n$ . Use the B-algorithm to synthesize a simple series-parallel circuit that computes  $115/512$  using only  $1/2$ -pswitches; draw the circuit. Draw a circuit that computes  $397/512$  using  $1/2$ -pswitches. How are the two circuits you synthesized related? Show your work.
- (b) What is the probability that the following series-parallel circuit is closed? Show your work. Can you increase the probability that the following circuit is closed by exchanging the probabilities of two pswitches? Show your work.



- (c) What is the probability that the following non-series-parallel circuits are closed? Assume that each pswitch is closed with probability  $1/2$ . Show your work.
- i.



ii.



## 2. Linear Threshold (LT) Gates (*Collaboration is allowed on this Problem*)

A linear threshold gate computes a Boolean function  $f(X)$  as follows,

$$f(X) = \text{sgn}(F(X)) = \text{sgn}\left(w_0 + \sum_{i=1}^n w_i x_i\right)$$

where,

$$\text{sgn}(F(X)) = \begin{cases} 0 & \text{if } F(X) < 0 \\ 1 & \text{if } F(X) \geq 0 \end{cases}$$

We say that a Boolean function is in the class  $LT_1$  if it can be computed by a single  $LT$  gate. An  $LT$  circuit is a circuit that consists of  $LT$  gates.

- (a) Recall (from HW#4) that the Complete Quadratic (CQ) function is the Boolean function that consists of the  $XOR$  of all the  $\binom{n}{2}$  possible  $AND$ 's between pairs of inputs. Namely,

$$CQ(X) = (x_1 \cdot x_2) \oplus (x_1 \cdot x_3) \oplus \cdots \oplus (x_{n-1} \cdot x_n)$$

Draw a depth-2 *layered*  $LT$  circuit that computes  $CQ(X)$  with six inputs, namely,  $X \in \{0, 1\}^6$ . Please use as few  $LT$  gates as possible. We expect you to use the method (due to Muroga) that was described in class.

- (b) We generalize linear threshold functions to accept more than two possible values in their inputs. Assume that every input  $x_i$  to a threshold function comes from an alphabet of size  $q$ , namely,  $x_i \in \{0, 1, 2, \dots, q-1\}$  (the binary case is when  $q = 2$ , namely,  $x_i \in \{0, 1\}$ ). For this generalization, we define the functions  $OR_q$  and  $AND_q$  as follows:

$$OR_q(x_1, x_2, \dots, x_n) = \begin{cases} 0 & \text{if } x_i = 0 \text{ for all } 1 \leq i \leq n \\ 1 & \text{otherwise} \end{cases}$$

and

$$AND_q(x_1, x_2, \dots, x_n) = \begin{cases} 1 & \text{if } x_i \neq 0 \text{ for all } 1 \leq i \leq n \\ 0 & \text{otherwise} \end{cases}$$

- i. For which  $q \geq 2$  is  $OR_q \in LT_1$ ? Prove your claim.
- ii. For which  $q \geq 2$  is  $AND_q \in LT_1$ ? Prove your claim.